

Concert Investigation of Novel Deterministic Interleaver for OFDM-IDMA System

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Abstract

In recent days the expectations of wireless communication systems have to configure new technologies that support high capacity to match the increase in demands on wireless service. A new hybrid scheme namely Orthogonal Frequency Division Multiplexing Interleave division Multiple Access (OFDM-IDMA) is a hopeful candidate of future wireless communication systems. In OFDM-IDMA the users are separated by having a unique interleaver pattern. Thus design of interleaver plays a prominent role in user separation. In this paper, new methodology and design analysis of novel deterministic interleaver is discussed. Two different deterministic interleavers are proposed and their collision probability is studied using correlation analysis. We also analyze the multiple access interference (MAI) performance of proposed interleavers using peak correlation analysis. The new deterministic interleavers are namely modified circular shifting interleaver and clockwise interleaver. The BER Vs E_b/N_0 analysis is carried out for the proposed interleavers and their performance is studied.

Keywords: Interleaver, OFDM, IDMA, OFDM-IDMA, Deterministic interleaver, circular shift interleaver, clockwise interleaver.

I. Introduction

Wireless communication system has grown tremendously with remarkable research effort and significant developments. Intersymbol interference (ISI) due to multipath, cross antenna interference caused due to multiple antenna transmission and multi user interference (MUI) occurred due to multiple user signal synchronized transmission are the major vulnerabilities of the wireless communication system. For Future applications like 4-G or next generation broadband wireless communication, a system combined with multicarrier modulation which is very effective in combating severe multipath interference and providing multiple access simultaneously is mandatory. Recent development shows the advantage of combining orthogonal frequency-division multiplexing (OFDM) and interleave division multiple-access (IDMA) in an OFDM-IDMA scheme. OFDM suppresses the inter-symbol interference (ISI) induced by multi-path channels, and the signals of different users (MUI) are avoided by user-specific interleavers based on the IDMA principle.

In OFDM-IDMA, IDMA (Interleave Division Multiple Access) is a novel technology that replaces the drawbacks of existing Code Division Multiple Access (CDMA) technique. In CDMA, interleaver is used for coding gain while in IDMA; they are engaged for user separation. The basic principle of IDMA is to use specific interleaver to differentiate the user. In order to immunize noise and Multiple Access Interference (MAI) at the receivers, selection of interleavers which are weakly correlated between different users is an important issue. The circumstance for IDMA to be successfully executed is that the interleaver between the transmitter and receiver should be identical. The desirable characteristics of IDMA interleaver codes includes (i) availability of large number of codes (ii) impulsive auto-correlation function (iii) zero cross-correlation values (iv) randomness (v) ease of generation (vi) Less hardware requirements (vii) High speed.

The stipulation for OFDM-IDMA to be successfully employed is that the transmitter and receiver concur upon the equivalent interleaver. Earlier OFDM-IDMA depended only on the random interleaver pattern. In the random interleavers; the entire interleaver matrix has to be transmitted to the receiver, which can be very costly. When the interleavers were generated randomly and self-regulating, the cross correlation between different interleavers is quite weak. The system requires more memory to hoard these interleaving patterns, because all of them will be used at the receiver for signal detection. Thus bandwidth must be reserved to modernize the interleaving information for OFDM-IDMA systems. When user number increases in the system,

this issue becomes worse. Meanwhile, the interleaving procedure should be fast enough to achieve high speed data transmission for the OFDM-IDMA systems.. In order to immunize noise and MAI for the receivers in OFDM-IDMA systems, it is necessary to choose good interleavers which are weakly correlated between different users. Our aspiration is to construct a methodical interleaver pattern for IDMA such that it performs as well as random interleavers and satisfy the design criteria: • The interleaver should be easy to stipulate and spawn, i.e., the transmitter and receiver requires a small size of bits to be in agreement upon an interleaver, and then generate it. • The interleavers should not have a collision.

The random interleaver introduced by Li Ping [1] can randomize the MAI at the receiver end efficiently. But user side and base station (BS) have to store a large number of interleaving pattern information. Thus, the memory and bandwidth requirements are ridiculous for systems. The nested interleaver by Hao Wu [2], consists of a basic interleaver, with which a series of interleavers can be obtained with recursion method. Comparing with the random interleavers, the memory and bandwidth resources consumption will be greatly reduced. The orthogonal interleavers and the pseudo random interleavers were introduced by Pupeza [3]. The design of orthogonal interleavers was based on the PN sequences. The pseudo random interleavers for different users depend on the primitive polynomials. Thus the specific user needs to store their primitive polynomial. Obviously, this method is better than the nested and the random interleavers. When user number becomes large, it is difficult to find enough primitive polynomials for pseudo random interleavers, and systems have to endure generation time for nested interleavers. Thus, pseudo random interleavers and nested interleavers are not suitable for large scale OFDM-IDMA system. The circular shifting interleaver [4] does a good job of permuting weight-2 input sequences with low codeword weights into weight-2 input sequences with high codeword weights. However, because of the regularity inherent in this type of interleaver, it may be difficult to permute higher weight (weight>2) input sequences with low codeword weights into other input sequences with high codeword weights.

In this paper, we propose the use of the deterministic interleaver. The receiver will by design generate the interleaver based on the user index k . This paper, deals with the design and peak correlation analysis of a deterministic chip-level interleaver. Neither the entire random interleaver matrix, nor the power interleaver needs to be transmitted. In the subsections the system model of OFDM-IDMA is explained in section II. It is followed by the design of deterministic interleaver in Section III. The MAI analysis is mentioned in section IV. Peak correlation and correlation analysis is discussed in Section V and VI. Simulation analysis is reported in Section VII which is followed by a conclusion in Section VIII.

II. System model

OFDM-IDMA [5] inherits most of the merits of OFDM and IDMA system. The key advantage of OFDM-IDMA is that multi user detection (MUD) can be realized efficiently with complexity per user independent of the channel length and the number of users. The transmitter/receiver structure of an OFDM-IDMA system with K users is shown in Figure 1.

This section briefs about the functioning OFDM-IDMA transmitter and receiver. OFDM-IDMA consists of a base station and number of users who transmits the data. Let the available bandwidth be divided among say k users. Thus the sub carriers are also k . The information bit from a single user is encoded by using low Forward Error Correcting Code [FEC] and a spreader. A spreader is simply a repetition code it consists of +1 and -1 alternatively e.g. {+1, -1, +1, -1, ...}. Then the coded data is interleaved using a specific interleaver. The interleaved data is mapped to a sequence using a modulator. Here a BPSK modulator is preferred. The sequence is modulated onto sub carriers by using an N point IFFT.

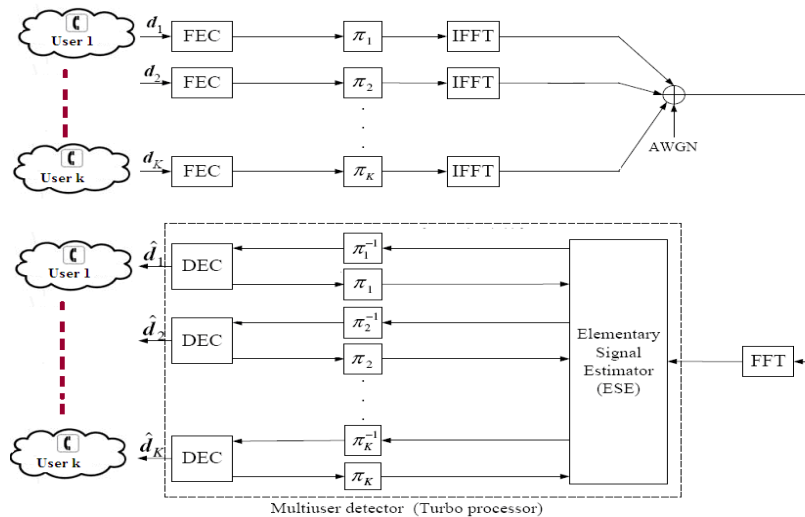


Fig 1 Modified model of OFDM-IDMA system is shown. Usage of modulator ,Cyclic prefix are not shown in the figure for simplicity.

The signal received at the base station is a superposition of all active users. The received signal is passed through OFDM demodulation network .The demodulation network comprises of FFT and removal of cyclic prefix. Then the signal is passed through iterative IDMA. Iterative IDMA comprises of Elementary signal estimator and Decoders. The process that is executed in iterative IDMA is explained as follows, the signal from OFDM demodulator is passed through Elementary Signal Estimator [ESE]. The ESE $e_{ESE}(X_k(j))$ process generates posterior Log Likelihood ratio (LLRs) of the given signal (symbol).

$$e_{ESE}(X_k(j)) = \log \left(\frac{p(X_k(j) = +1|r)}{p(X_k(j) = -1|r)} \right) \nu_{k,j} \tag{1}$$

$$e_{ESE}(X_k(j)) = 2h_k \left(\frac{r(j) - E(\zeta_k(j))}{\text{var}(\zeta_k(j))} \right) \tag{2}$$

Where, $E(X_k(j))$ is the mean and $\text{Var}(X_k(j))$ is the variance of the transmitted bits.

$$E(X_k(j)) = \tanh(e_{DEC}(X_k(j)))/2 \tag{3}$$

$$\text{Var}(X_k(j)) = 1 - (E(X_k(j)))^2 \tag{4}$$

$e_{ESE}(X_k(j))$ is sent to the APP-DEC after de-interleaving. And the e_{dec} are generated by the APP-DEC is sent to the ESE through the interleaving, the process is iterated several times next iteration. Once the iteration is over the decoder produces hard decision to estimate the transmitted bits.

III. Design Analysis of Deterministic Interleaver

In this paper we have introduced two new deterministic interleaver namely modified circular shifting interleaver and clockwise interleaver. Deterministic interleavers [6] are constructed using a mathematical expression. The receiver automatically generates the interleaver based on the user index k. It is generated easily and time taken for generation is feasible. This aims at reducing the bandwidth and the memory requirement of the systems which employ random interleavers.

a. Modified Circular shift Interleaver.

(i) Methodology

This interleaver is acquired from circular shift interleaver [4] which was defined by the equation $P(i) = (a(i) + s) \bmod L$, Where a is a prime number, I is the index and s is the offset and $a < L$ and $s < 1$. Modified circular shift interleaver [6] is specified as

$$K = \prod [(P * Q) \text{ MOD } N] \tag{5}$$

The deinterleaver is obtained by using the symmetry property of modulus arithmetic

$$Q = \prod^{-1} [(K \text{ MOD } N) / P] \tag{6}$$

\prod represents the interleaver and \prod^{-1} denotes the deinterleaver. Generally \prod and \prod^{-1} are self inverse to each other. P denotes any prime number and N denotes the size of the coded data. Also Q and K denotes the bit position. Q resembles the existing bit position of the coded data and K denotes the new interleaver position. Normally the prime number P chosen should not be a factor of N. In that case interleaving pattern is not obtained. Depending upon the number of user to be allotted distinctive prime number is chosen for each user.

(ii) Design Procedure

Step 1: The length of the interleaver sequence is N. Let the bit position q, of the interleaver be 0 to N-1.

Step 2: Allot a unique prime number to each user. A prime number chosen should not be a factor of N.

Step 3: The interleaved bit position is obtained using equation 1.

Step 4: The deinterleaved pattern is obtained using equation 2.

The flow chart shown in Figure 2 briefs the procedure explained above in design procedure

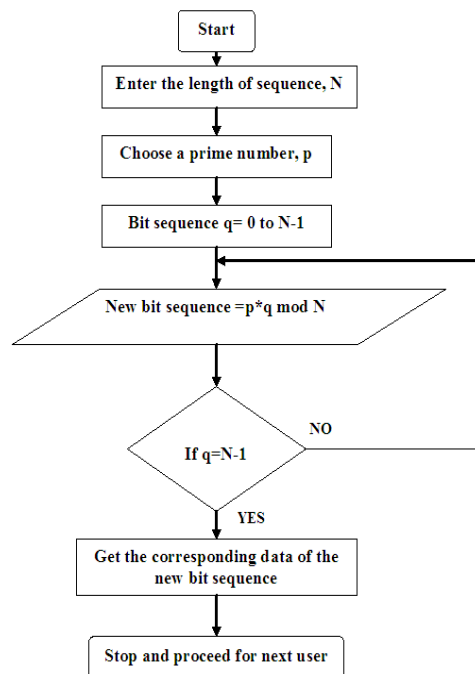


Fig 2. Flow chart of modified circular shift interleaver using (P*Q) mod N is depicted.

(iii) Example

Consider N to be 8 bits such that {0, 1, 2, 3, 4, 5, 6, 7}. Let the prime number be 3 then the new location of bit will be as follows let

- 0====> (0*3) mod 8= 0
- 1====> (1*3) mod 8 =3
- 2====> (2*3) mod 8= 6
- 3====> (3*3) mod 8= 1
- 4====> (4*3) mod 8= 4
- 5====> (5*3) mod 8= 7
- 6====> (6*3) mod 8= 2
- 7====> (7*3) mod 8= 5

Now, the new order of bits will be { 0,3,6,1,4,7,2,5}.

b. Clockwise Interleaver.

(i) Methodology

Another deterministic interleaver mentioned here is different from the above one in the type of seed used. Instead of prime number the user index itself is used as the seed.

$$[] = (X+k) \text{ mod } L \tag{7}$$

X indicates the user number and k indicates the bit position with L has the length of the interleaver.

$$[]^{-1} = (X+m) \text{ mod } L \tag{8}$$

The de-interleaver is obtained by using the symmetry property of modulus arithmetic

(ii) Design procedure

Step 1: The length of the interleaver sequence is L.Let the bit position k, of the interleaver be 0 to L-1..

Step 2: The user index is assigned as seed. Assign the user number itself as X

Step 3: The interleaved bit position is obtained by equation (7).

Step 4: The de-interleaved bit position is obtained by equation (8).

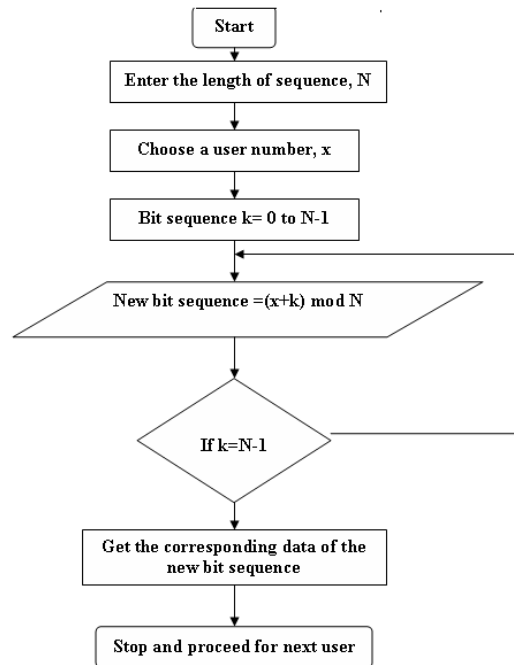


Fig 3. Flow chart of clockwise interleaver using (X+K) mod N is depicted.

(iii) Example:

Consider the user number to be 3 and the consecutive bits position K be {0, 1, 2, 3, 4-----L-1}. Assume L to be 8 then the new location of the bit becomes,

- 0====> (0+3) mod 8= 3
- 1====> (1+3) mod 8 =4
- 2====> (2+3) mod 8= 5
- 3====> (3+3) mod 8= 6
- 4====> (4+3) mod 8= 7
- 5====> (5+3) mod 8= 0
- 6====> (6+3) mod 8= 1
- 7====> (7+3) mod 8= 2

Now, the new order of bits will be { 3,4,5,6,7,0,1,2}.

The benefit of clockwise interleaver is that number of user allotted depends on the size of the interleaver.The flow chart shown in Figure 3 briefs the procedure explained above in design procedure.

IV. Analysis of Multiple Access Interference (MAI)

Assume that there are only two users in the system and noiseless channel is used. The data of user1 is encoded by C and permuted by Π_1 . At the receiver end the received signal is $C_1 * \Pi_1 + C_2 * \Pi_2$. The received signal is de-interleaved by Π_1^{-1} . Now the input information of decoder is $C_1 + C_2 * \Pi_2 * \Pi_1^{-1}$. If C_1 and $C_2 * \Pi_2 * \Pi_1^{-1}$ are orthogonal to each other, than the interference from user2 can be eliminated. To suppress the MAI (Multiple Access Interference) the cross correlation between them need to be zero.

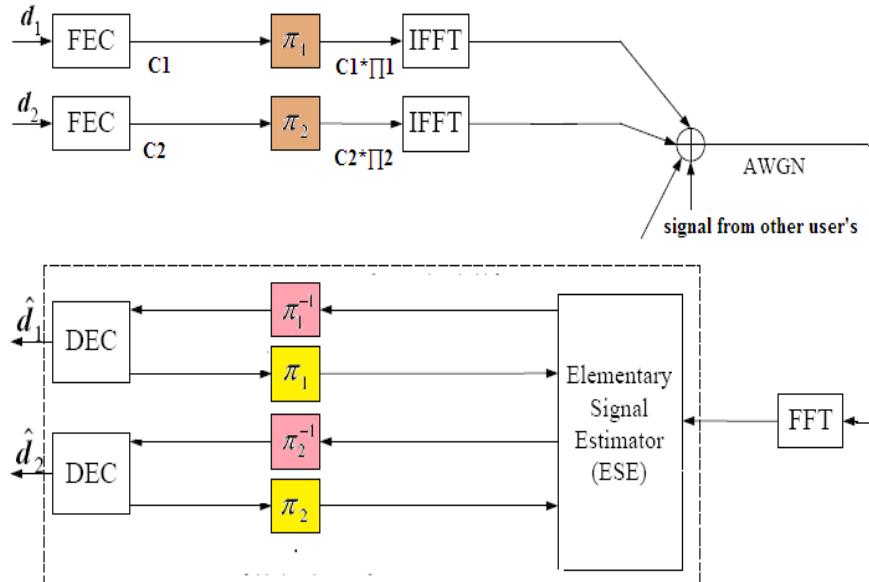


Fig 4 .System model of OFDM-IDMA with two users is depicted with hidden modular and cyclic prefix blocks.

$$C (C_1 * \Pi_1, C_2 * \Pi_2) = 0 \tag{9}$$

So the orthogonal interleavers can be defined as following. For any two coded words, C_i and C_j are permuted by two different interleavers Π_1 and Π_2 respectively, if their cross correlation is zero

$$C (C_i * \Pi_1, C_j * \Pi_2) = 0 \tag{10}$$

The orthogonal attribute of OFDM-IDMA systems has bond with coded data pattern and interleaving scheme. Thus, it is complicated to achieve orthogonal interleavers in the practical OFDM-IDMA systems. Generally, weakly correlated interleavers can be used in the OFDM-IDMA system to obtain good performance. Since correlation of different interleavers is affected by the coded data pattern, we use Peak Correlation [3] to analyze the correlation issues for simplification in this paper.

V. Peak Correlation

(1) Peak Auto-Correlation

Auto-correlation is a measure of the similarity between a code $C(t)$ and its time shifted replica [7]. Mathematically defined as

$$\psi(\tau) = \int_{-\infty}^{\infty} c(t) * c(t - \tau) dt \tag{11}$$

Ideally, this auto-correlation function (ACF) should be spontaneous i.e. peak value at zero time shift and zero values at all other time-shifts (i.e. side-lobes). This is required at the receiver side for proper harmonization and to discriminate the desired user from other users producing MAI.

(2) Peak Cross-Correlation

Cross-correlation [7] is the measure of similarity between two different code sequences $C_1(t)$ and $C_2(t)$. Mathematically, it is defined as:

$$\psi(\tau) = \int_{-\infty}^{\infty} c_1(t) * c_2(t - \tau) dt \quad (12)$$

Cross-correlation function (CCF) designates the correlation between the desired and undesired code sequences at the receiver. Therefore, in order to eliminate the effect of multiple access interference at the receiver, the cross-correlation value is obliged to zero at all time shifts.

VI. Correlation Analysis

When two interleaver sequences are strongly related then the correlation between them is '1'. When two interleavers are orthogonal to each other, then the correlation between them is "0". The correlation of any two sequences lies between -1 to +1. If the correlation value approaches "0" then these sequences do not match each other. For orthogonality condition the value should be zero. The correlation for different user specific interleaver [6] may be defined as Let data1 and data2 be the two data sequences of two different user, say user1 and user2. Then find XOR of user 1 and user 2 as data1 XOR data2. Now calculate the amount of 1's after XORing the both user data as X and the amount of 0's after XORing the both user data as Y. Now correlation is given as

$$\text{Correlation value} = (X - Y) / (X + Y) \quad (13)$$

The Peak Correlation numerical results of modified circular shifting interleavers and clockwise interleaver are calculated. For comparison, the Peak Correlation performances of orthogonal interleaver, modified circular shift interleaver and clockwise interleaver are also studied.

VII. Simulation Result's of Proposed Deterministic Interleaver

(1) Peak Correlation Analysis:

The Peak Correlation numerical results are calculated for a sample of interleaver size 16 for five users. The prime numbers allotted for each user is 3, 5, 7, 11 and 13 respectively. The data's observed are given in the tables below as follows

Table1: Peak correlation of an ideal orthogonal Interleaver

USER	1	2	3	4	5
1	16	0	0	0	0
2	0	16	0	0	0
3	0	0	16	0	0
4	0	0	0	16	0
5	0	0	0	0	16

Table2: peak correlation of modified circular interleaver

USER	1	2	3	4	5
1	16	6	6	8	10
2	6	16	6	6	8
3	6	6	16	6	6
4	8	6	6	16	6
5	10	8	6	6	16

Table3: peak correlation of clockwise interleaver

USER	1	2	3	4	5
1	16	6	8	2	6
2	6	16	10	6	2
3	8	10	16	8	10
4	2	6	8	16	6
5	6	2	10	6	16

The diagonal elements in the above tables are auto-correlations and the other elements are the cross correlation between different users. From these tables above, we can see that the correlation values of modified shifting and clockwise interleavers are similar. The auto correlation has peak value at zero shifts. The maximum weight of the data is 10(1111010000111110).The peak cross correlation value lies between 1 and 10. When the length of the data is increased the cross correlation reduces close to zero.

(2) Correlation Analysis

(i) Clockwise interleaver

For clockwise interleaver thirty users were considered with a data length of eight .The correlation is as follows

Columns 1 through 8

0.06 0.33 -0.2 -0.2 0.06 0.33
-0.2 -0.2

Columns 9 through 16

0.33 0.06 -0.2 -0.2 0.33 0.06
0.06 0.33

Columns 17 through 24

-0.2 -0.2 0.33 0.06 -0.2 -0.2
0.33 0.06

Columns 25 through 29

-0.2 -0.2 0.33 0.06 -0.2 -0.25

User specific cross correlation is 0.27. The above values are simulated and plotted in figure 5 using matlab.

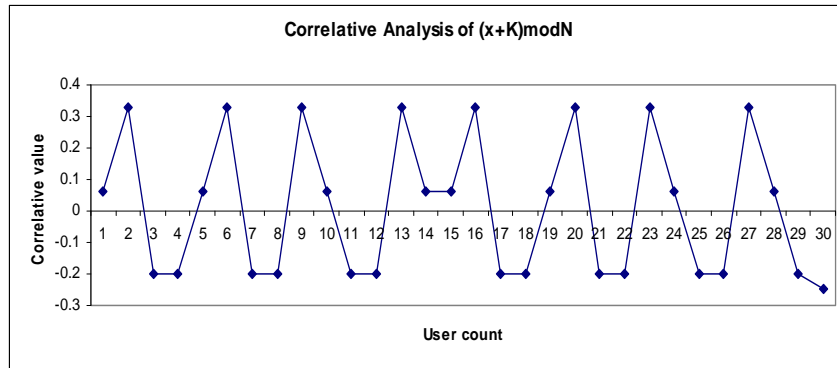


Fig 5 Correlation Analysis of (x+k) mod N interleaver with thirty user

(ii) Modified circular shift interleaver

For circular shifting interleaver thirty users were considered with a data length of eight. The correlation is as follows

Columns 1 through 8

0.25 0.25 0.5 0.25 -0.25 -0.25 -0.25 -0.5

Columns 9 through 16

-0.75 -0.25 0 0.5 -0.25 -0.25 0.25 0.25

Columns 17 through 24

0.5 0.25 -0.25 -0.25 -0.25 -0.5
-0.75 -0.25

Columns 25 through 29

0 0 0.5 -0.25 -0.25 0.75

User specific cross correlation is -0.5. The above values are simulated and plotted in figure 6 using matlab.

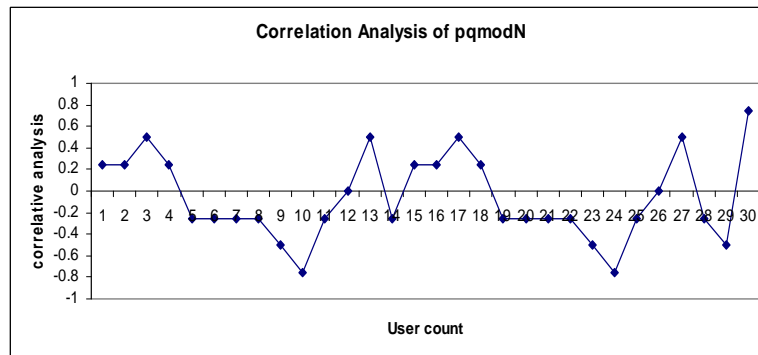


Fig 6. Correlation Analysis of (p*q) mod N for thirty user's.

(iii) Comparison chart

The comparison between modified circular shifting interleaver and clockwise interleaver is showed in figure 8. we can examine that clockwise interleaver has perfect correlation between the interleavers than the modified circular shifting interleaver.

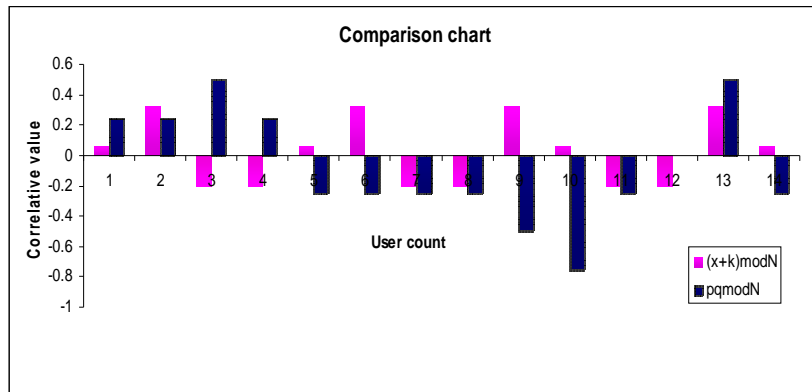


Fig 7. The chart enhances the comparison between the performance of both proposed interleaver with thirty user's

(3) BER Vs Eb/N0 analysis

The performance analysis of modified circular shift interleaver is carried out using basic BPSK modulator. The number of bits per symbol considered is 52. The analysis is carried out for a maximum of 35 user. Number of symbols considered is 10,000. Size of Fast Fourier Transform is fifty four. The Forward error correcting code used is 1/2 convolutional encoder. The Bit error rate is reduced when the number of user introduced.

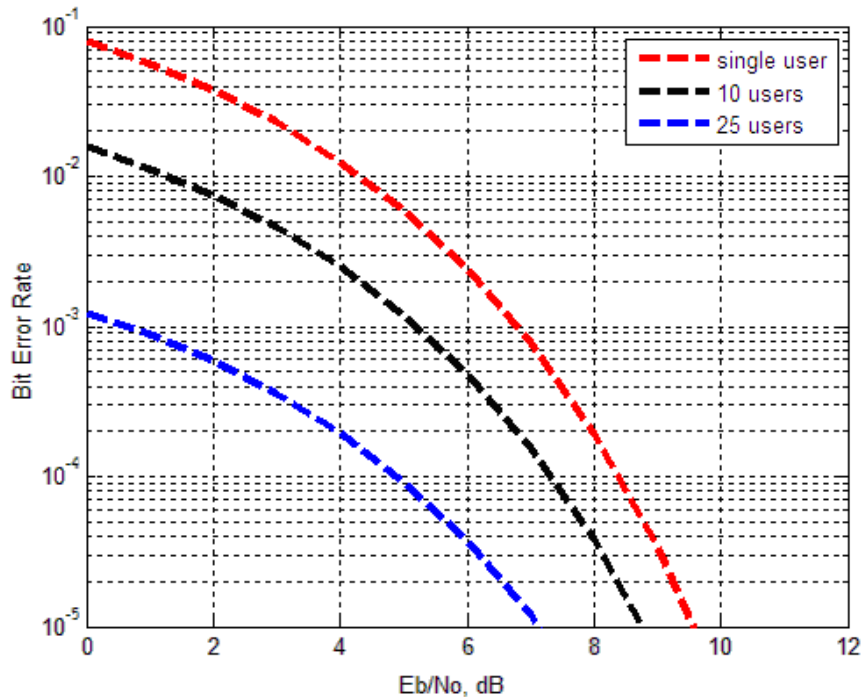


Fig 8. Performance Analysis of modified circular shift interleaver with thirty users and the modulation scheme used is BPSK.

The performance analysis of clockwise interleaver is carried out using QAM modulator. The number of bits per symbol considered is 52. The analysis is carried out for a maximum of thirty users. The Forward error correcting code used is 1/2 convolutional encoder. The Bit error rate is reduced when the number of user is increased .

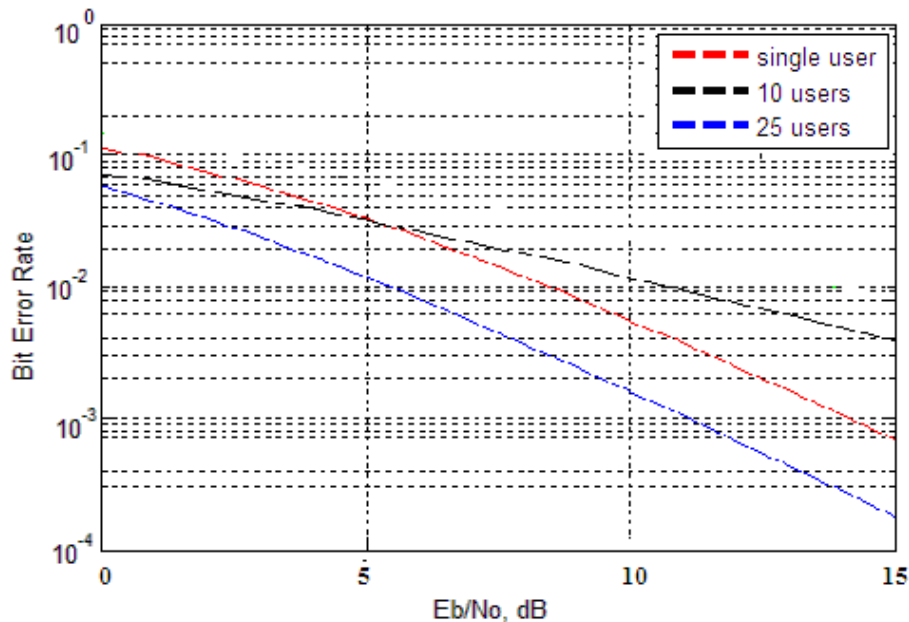


Fig 9. Performance Analysis of clockwise interleaver with thirty users and the modulation scheme used is QAM.

VIII. Inference

The correlation value of random interleaver increases with the no. of users'. The deterministic interleaver indicates a correlation value less than one. The clockwise interleaver has a user specific cross correlation 0.27 which is nearly zero. This indicates that the collision between two interleavers is nearly zero. In the circular shifting interleaver the user specific cross correlation value is -0.5. The correlation value between -1 to 0 is acceptable. The performance of modified circular shift interleaver is better when compared to clockwise interleaver. In the future ventures more study analysis has to be executed with common higher modulation scheme. In future more deterministic interleavers are to be recognized and their performance should be studied.

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